

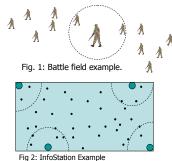
# Cooperative Cache Based Data Access in Ad Hoc Networks



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### Motivation

- Data access issues in ad hoc networks: Battle field:
  - Long query delay due to multi-hop links
  - Low data availability due to node/link failure
- Cooperative caching schemes
  - Share and coordinate cached data.
  - Reduce guery delay, bandwidth/power consumption.
- - Officer may have a powerful data center. Soldiers need to access the data center to get geographic info, enemy info, and new commands. Neighboring soldiers share these info.
- InfoStations
  - Infostations are deployed in cities to provide info, such as maps, attractive sites, or restaurant info to mobile users. Users may relay for each other to serve those not directly covered by the infostations.



Green node: infostation; Black node: mobile node

### Data Access Framework

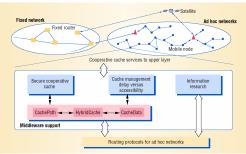


Fig. 3: Cooperative cache-based data access framework

- The cooperative cache-based data access framework stays on top of the routing protocol.
- It relies on several components
  - Secure cooperative caching,
  - Cache Management,
  - Information search

to provide services to upper layer.

### System Model



Fig. 4: Ad hoc networks. Node  $N_{11}$  is a data center and the blue nodes are router nodes. Node  $N_1$  is a cluster header surrounded by cluster member nodes.

### The Cache Protocols

#### The CacheData Scheme

- Router nodes caches frequently accessed passing-by data to serve future requests.
  - An example (using Fig. 4)
    - Suppose both N<sub>6</sub> and N<sub>7</sub> requested data item  $d_i$  through  $N_5$ .
    - $N_5$  might think that  $d_i$  is popular and cache it locally.
    - N<sub>5</sub> can then serve N<sub>4</sub>'s request directly.

#### The CachePath Scheme

- Cache the node id that requests the data when a data item passes by.
  - An example (using Fig. 4)
    - Suppose N<sub>1</sub> requests d<sub>i</sub> from N<sub>11</sub>.
    - When N<sub>3</sub> forwards d<sub>i</sub> to N<sub>1</sub>, it caches the id pair of  $N_1$  and  $d_i$ : (i, 1).
    - Future requests for d<sub>i</sub> from N<sub>2</sub> can be redirected to N₁ to reduce query delay and resource consumption.

#### HybridCache- a Hybrid Scheme



Fig. 5: HybridCache: take advantage of both CacheData and

## Cache Management

- Cache replacement policies
  - To determine whether to replace d<sub>i</sub> node Ni
    - δ: Distance of Ni to the data center or a caching node of di, small  $\delta$  value is preferred for removal.
    - ζ: last update time of δ, obsolete data is preferred.
    - ullet Combinations of  $\delta$  and  $\zeta$  can affects the tradeoffs between query delay and data accessibility.
- Cache admission control
  - Following similar rules used in cache replacement

### **Experimental Result**



### Conclusion

- A cooperative cache-based data access framework
  - Let mobile nodes cache the data or path to the data to reduce the query delay and improve data accessibility
  - Provide secure cooperative cache, cache management, and information search services to upper layer.